

A Method for Cross-layer Analysis of Transmit Buffer Delays in Message Index Domain

WoNeCa

Sami Akın and Markus Fidler



Institute of Communications Technology
Leibniz Universität Hannover

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Outline

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Motivation

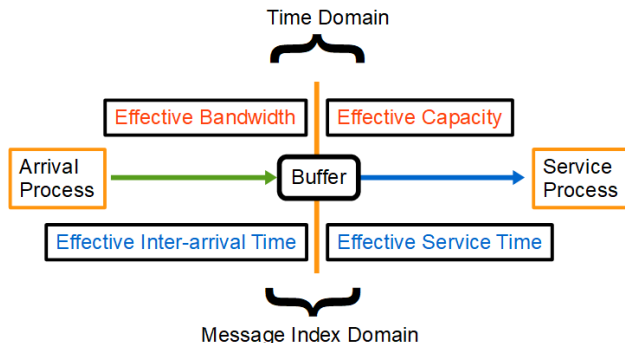
- Need for faster and more reliable data transmission techniques
- Effective bandwidth and capacity concepts
- An approximation for the steady-state buffer overflow probability

- Delay violation probability of first-come first-served systems when either the data arrival or service rate is constant
- Difficult to obtain effective bandwidth and capacity in certain scenarios without resorting to numerical techniques or particular assumptions

Contributions

- Analysis in the message index domain
- Waiting time (buffering delay) characterization

- Waiting time violation probability



Assuming that the buffer is unbounded (very big), the backlog (queue length) at time instant t is

$$q(t) = [q(t-1) + a_q(t) - s_q(t)]^+$$

$a_q(t)$: arrival rate, $s_q(t)$: service rate

$$[\cdot]^+ = \max\{0, \cdot\}$$

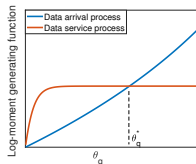
$q(\infty)$ will converge in distribution to a finite random variable when both $a_q(t)$ and $s_q(t)$ are stationary and ergodic, and when $\mathbb{E}\{a_q(t)\} < \mathbb{E}\{s_q(t)\}$

When $a_q(t)$ and $s_q(t)$ are independent of each other, there exists a unique θ_q^* such that

$$\lim_{t \rightarrow \infty} \frac{\log \mathbb{E} \left\{ e^{\theta_q^* \sum_{\tau=1}^t a_q(\tau)} \right\}}{t} = - \lim_{t \rightarrow \infty} \frac{\log \mathbb{E} \left\{ e^{-\theta_q^* \sum_{\tau=1}^t s_q(\tau)} \right\}}{t}$$

and

$$\theta_q^* = - \lim_{q_{th} \rightarrow \infty} \frac{\log \Pr\{q(\infty) \geq q_{th}\}}{q_{th}}$$



For large q_{th} , $\Pr\{q(\infty) \geq q_{th}\} \approx e^{-\theta_q^* q_{th}}$

For $\theta_q > 0$

$$\text{Effective bandwidth: } \bar{a}_q(\theta_q) = \lim_{t \rightarrow \infty} \frac{\log \mathbb{E} \left\{ e^{\theta_q \sum_{\tau=1}^t a_q(\tau)} \right\}}{\theta_q t}$$

$$\text{Effective capacity: } \bar{s}_q(-\theta_q) = \lim_{t \rightarrow \infty} \frac{\log \mathbb{E} \left\{ e^{-\theta_q \sum_{\tau=1}^t s_q(\tau)} \right\}}{-\theta_q t}$$

Given a constant service or arrival rate;

$$\Pr\{d(\infty) \geq d_{\text{th}}\} \approx \varepsilon e^{-\theta_q \bar{a}_q(\theta_q) d_{\text{th}}} \quad \text{or} \quad \varepsilon e^{-\theta_q \bar{s}_q(-\theta_q) d_{\text{th}}}$$

where

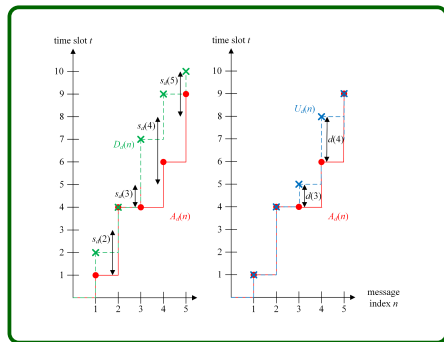
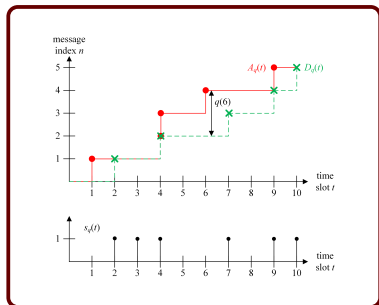
$$d_{\text{th}} = \frac{q_{\text{th}}}{\bar{a}_q(\theta_q)} \quad \text{or} \quad \frac{q_{\text{th}}}{\bar{s}_q(-\theta_q)}, \quad \text{and} \quad d(\infty) = \frac{q(\infty)}{\bar{a}_q(\theta_q)} \quad \text{or} \quad \frac{q(\infty)}{\bar{s}_q(-\theta_q)}$$

Objectives and approach

- $\Pr\{d(\infty) \geq d_{\text{th}}\}$ when both arrival and service processes are stochastic
- Communication scenarios where effective bandwidth and effective capacity are not straightforward

The pseudo-inverse functions of cumulative arrival and service processes: the number of messages and the number of time instants are positioned on the x-axis and y-axis, respectively

Objectives and approach



- $A_d(t)$: Cumulative arrival process
- $D_d(t)$: Cumulative departure process

- $A_d(n)$: Message n arrives
- $D_d(n)$: Message n departs
- $U_d(n)$: Message n starts service

Waiting time characterization

The waiting time:

$$d(n) = [d(n-1) + s_d(n) - a_d(n)]^+ \text{ for } n \geq 1$$

$d(n)$: waiting time of message n

$a_d(n)$: Inter-arrival time between messages $n-1$ and n

$s_d(n)$: Service time of message $n-1$

Time message n starts service:

$$U_d(n) = \max \{A_d(n), U_d(n-1) + s_d(n)\}$$

$A_d(n) = \sum_{i=1}^n a_d(i)$ is the time message n arrives at the buffer

Steady-State Waiting Time Analysis

$d(\infty)$ will converge in distribution to a finite random variable when both $s_d(n)$ and $a_d(n)$ are stationary and ergodic, and when $\mathbb{E}\{s_d(n)\} < \mathbb{E}\{a_d(n)\}$

The service process is work-conserving, and the arrival and service processes are independent of each other;

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log \mathbb{E} \left\{ e^{\theta_d^* \sum_{i=1}^n s_d(i)} \right\} = - \lim_{n \rightarrow \infty} \frac{1}{n} \log \mathbb{E} \left\{ e^{-\theta_d^* \sum_{i=1}^n a_d(i)} \right\}$$

where

$$\theta_d^* = - \lim_{d_{\text{th}} \rightarrow \infty} \frac{\log \Pr\{d(\infty) \geq d_{\text{th}}\}}{d_{\text{th}}}$$

and large d_{th}

$$\Pr\{d(\infty) \geq d_{\text{th}}\} \approx e^{-\theta_d^* d_{\text{th}}}$$

Effective service time

For $\theta_d > 0$

$$\bar{s}_d(\theta_d) = \lim_{n \rightarrow \infty} \frac{1}{n\theta_d} \log \mathbb{E} \left\{ e^{\theta_d \sum_{i=1}^n s_d(i)} \right\}$$

When the service time samples are independent and identically distributed;

$$\bar{s}_d(\theta_d) = \frac{1}{\theta_d} \log \mathbb{E} \left\{ e^{\theta_d s_d(n)} \right\}$$

$$\lim_{\theta_d \rightarrow 0} \bar{s}_d(\theta_d) = \mathbb{E} \{ s_d(n) \} \quad \text{and} \quad \lim_{\theta_d \rightarrow \infty} \bar{s}_d(\theta_d) = \max \{ s_d(n) \}$$

Effective inter-arrival time

For $\theta_d > 0$

$$\bar{a}_d(-\theta_d) = \lim_{n \rightarrow \infty} \frac{-1}{n\theta_d} \log \mathbb{E} \left\{ e^{-\theta_d \sum_{i=1}^n a_d(i)} \right\}$$

When the inter-arrival time samples are independent and identically distributed;

$$\bar{a}_d(-\theta_d) = -\frac{1}{\theta_d} \log \mathbb{E} \left\{ e^{-\theta_d a_d(n)} \right\}$$

$$\lim_{\theta_d \rightarrow 0} \bar{a}_d(-\theta_d) = \mathbb{E} \{ a_d(n) \} \quad \text{and} \quad \lim_{\theta_d \rightarrow \infty} \bar{a}_d(-\theta_d) = \min \{ a_d(n) \}$$

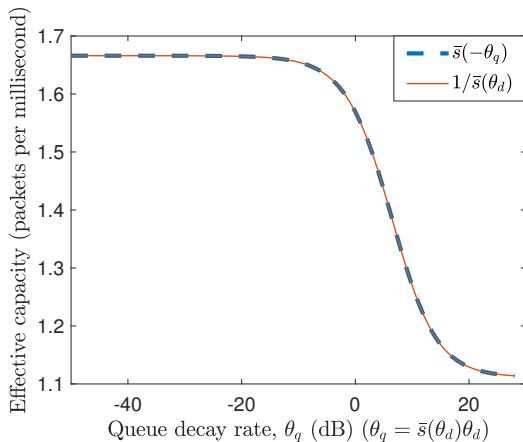
Message index domain \iff Time domain

Given that the effective service time of a data service process is $\bar{s}_d(\theta_d)$ for $\theta_d > 0$ and the effective capacity of the same service process is $\bar{s}_q(-\theta_q)$ for $\theta_q > 0$, there exists a reciprocal relation between $\bar{s}_d(\theta_d)$ and $\bar{s}_q(-\theta_q)$ such that

$$\bar{s}_d(\theta_d)\bar{s}_q(-\theta_q) = 1 \quad \text{and} \quad \theta_q = \bar{s}_d(\theta_d)\theta_d$$

Given that the effective inter-arrival time of a data arrival process is $\bar{a}_d(-\theta_d)$ for $\theta_d > 0$ and the effective bandwidth of the same arrival process is $\bar{a}_q(\theta_q)$, there exists a reciprocal relation between $\bar{a}_d(-\theta_d)$ and $\bar{a}_q(\theta_q)$ such that

$$\bar{a}_d(-\theta_d)\bar{a}_q(\theta_q) = 1 \quad \text{and} \quad \theta_q = \bar{a}_d(-\theta_d)\theta_d$$



Thank You

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Questions