



Improving Output Bounds in the Stochastic Network Calculus Using Lyapunov's Inequality

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Current State of SNC reaching a Crossroads



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MGF-Calculus

Theorem

The violation probability of a given stochastic delay bound T at time t is bounded by

$$P(d(t) > T) \leq E\left[e^{\theta(A \otimes S)(t+T,t)}\right]$$

= $E\left[\max_{\substack{0 \leq i \leq t}} e^{\theta(A(i,t)-S(i,t+T))}\right]$
 $\stackrel{(\mathsf{UB})}{\leq} \sum_{i=0}^{t} E\left[e^{\theta(A(i,t)-S(i,t+T))}\right]$

Purpose of the union bound in the analysis: get rid of the max.

Union Bound

- does not need additional assumptions, independent of the distribution
- "is not going to be so bad if the variables (...) are rather uncorrelated" [Talagrand, 1996] or for Poisson arrivals [Ciucu, 2007]

Union Bound

- does not need additional assumptions, independent of the distribution
- "is not going to be so bad if the variables (...) are rather uncorrelated" [Talagrand, 1996] or for Poisson arrivals [Ciucu, 2007]
- is known to perform poorly for arrivals with correlated increments, such as Markov-Modulated On-Off traffic [Ciucu et al., 2013] [Beck, 2016]

Goal:

Mitigate the union bound's inaccuracy

$$\mathsf{E}\left[\max_{i=1,\ldots,n} e^{\theta X_i}\right] \leq \sum_{i=1}^n \mathsf{E}\left[e^{\theta X_i}\right]$$

but still obtain end-to-end delay bounds.

Lyapunov's Inequality

Proposition

Let $X \ge 0$ be in \mathcal{L}^{l} with $l \ge 1$. Then it holds that

$$\mathsf{E}[X] \leq \left(\mathsf{E}[X']\right)^{rac{1}{l}}.$$

Lyapunov's Inequality

Proposition

Let $X \ge 0$ be in \mathcal{L}^{I} with $I \ge 1$. Then it holds that

$$\mathsf{E}[X] \leq \left(\mathsf{E}\left[X'\right]\right)^{\frac{1}{l}}.$$

As I = 1 is feasible, this yields

$$\mathsf{E}[X] = \inf_{l \ge 1} \left\{ \left(\mathsf{E}[X^{l}] \right)^{\frac{1}{l}} \right\}.$$

$$\mathsf{P}(d(t) > T) \leq \mathsf{E}\left[e^{\theta(A \oslash S)(t+T,t)}\right] \\ = \mathsf{E}\left[e^{\theta\max_{0 \leq i \leq t}\left\{A(i,t) - S(i,t+T)\right\}}\right]$$

$$\begin{split} \mathsf{P}(d(t) > \mathcal{T}) &\leq \mathsf{E}\big[e^{\theta(A \oslash S)(t+\mathcal{T},t)}\big] \\ &= \mathsf{E}\big[e^{\theta\max_{0 \leq i \leq t}\{A(i,t) - S(i,t+\mathcal{T})\}}\big] \\ &\overset{(\mathsf{UB})}{\leq} \sum_{i=0}^{t} \mathsf{E}\big[e^{\theta(A(i,t) - S(i,t+\mathcal{T}))}\big] \end{split}$$

$$P(d(t) > T) \leq E[e^{\theta(A \otimes S)(t+T,t)}]$$

$$= E[e^{\theta \max_{0 \leq i \leq t} \{A(i,t) - S(i,t+T)\}}]$$

$$\stackrel{(\mathsf{LI})}{=} \inf_{l \geq 1} \left\{ \left(E[e^{l\theta \max_{0 \leq i \leq t} \{A(i,t) - S(i,t+T)\}}] \right)^{\frac{1}{l}} \right\}$$

$$P(d(t) > T) \leq E\left[e^{\theta(A \otimes S)(t+T,t)}\right]$$

$$= E\left[e^{\theta \max_{0 \leq i \leq t}\left\{A(i,t) - S(i,t+T)\right\}}\right]$$

$$\stackrel{(LI)}{=} \inf_{l \geq 1} \left\{ \left(E\left[e^{l\theta \max_{0 \leq i \leq t}\left\{A(i,t) - S(i,t+T)\right\}}\right]\right)^{\frac{1}{7}}\right\}$$

$$\stackrel{(UB)}{\leq} \inf_{l \geq 1} \left\{ \left(\sum_{i=0}^{t} E\left[e^{l\theta(A(i,t) - S(i,t+T))}\right]\right)^{\frac{1}{7}}\right\}$$

$$\begin{split} \mathsf{P}(d(t) > \mathcal{T}) &\leq \mathsf{E}\left[e^{\theta(A \otimes S)(t+\mathcal{T},t)}\right] \\ &= \mathsf{E}\left[e^{\theta\max_{0 \leq i \leq t}\left\{A(i,t) - S(i,t+\mathcal{T})\right\}}\right] \\ & \left(\mathsf{LI} \right) \\ & \left(\mathsf{II} \right$$

as l = 1 is feasible.

Significant Improvement, but wait a Minute...



Figure: Single node delay bound with MMOO arrivals and constant rate server

Lyapu-Fail?



Figure: Single node delay bound with MMOO arrivals and constant rate server

Numerical experiments revealed that no improvement at all is achieved for bounds < 1!</p>

Delay and Lyapunov Inequality are not meant to be together!

Proposition

Let T in the delay bound $P(d(t) > T) \leq \inf_{l \geq 1} \left\{ \left(\sum_{i=0}^{t} E[e^{l\theta(A(i,t) - S(i,t+T))}] \right)^{\frac{1}{l}} \right\} \text{ be large}$ enough such that the bound on the RHS is < 1. If l and θ are optimized, then the optimal l is 1, i.e., no improvement can be achieved with Lyapunov's inequality.





\Rightarrow Can it still be of any use?

- Cannot improve probability bounds, such as P(d(t) > T)
- Apply to MGF of output $\mathsf{E}\left[e^{\theta A'(s,t)}\right]$ instead

$$\mathsf{E}\Big[e^{\theta A'(s,t)}\Big] \leq \mathsf{E}\Big[e^{\theta \max_{0 \leq i \leq t} \{A(i,t) - S(i,s)\}}\Big]$$

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$$\overset{(\mathsf{UB})}{\leq} \sum_{i=0}^{t} \mathsf{E}\Big[e^{\theta(A(i,t) - S(i,s))}\Big]$$

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$$\stackrel{(\mathsf{L})}{=} \inf_{l \geq 1} \left\{ \left(\mathsf{E}\Big[e^{l\theta \max_{0 \leq i \leq t} \{A(i,t) - S(i,s)\}}\Big]\right)^{\frac{1}{l}} \right\}$$

$$E\left[e^{\theta A'(s,t)}\right] \leq E\left[e^{\theta \max_{0 \leq i \leq t}\left\{A(i,t) - S(i,s)\right\}}\right]$$

$$\stackrel{(LI)}{=} \inf_{l \geq 1} \left\{ \left(E\left[e^{l\theta \max_{0 \leq i \leq t}\left\{A(i,t) - S(i,s)\right\}}\right]\right)^{\frac{1}{7}}\right\}$$

$$\stackrel{(UB)}{\leq} \inf_{l \geq 1} \left\{ \left(\sum_{i=0}^{t} E\left[e^{l\theta(A(i,t) - S(i,s))}\right]\right)^{\frac{1}{7}}\right\}$$

$$\begin{split} \mathsf{E}\Big[e^{\theta A'(s,t)}\Big] &\leq \mathsf{E}\Big[e^{\theta \max_{0 \leq i \leq t}\{A(i,t) - S(i,s)\}}\Big] \\ \stackrel{(\mathsf{L}\mathsf{I})}{=} &\inf_{l \geq 1} \left\{ \left(\mathsf{E}\Big[e^{l\theta \max_{0 \leq i \leq t}\{A(i,t) - S(i,s)\}}\Big]\right)^{\frac{1}{l}} \right\} \\ \stackrel{(\mathsf{UB})}{\leq} &\inf_{l \geq 1} \left\{ \left(\sum_{i=0}^{t} \mathsf{E}\Big[e^{l\theta(A(i,t) - S(i,s))}\Big]\right)^{\frac{1}{l}} \right\} \\ \stackrel{\vdots}{\leq} &\sum_{i=0}^{t} \mathsf{E}\Big[e^{\theta(A(i,t) - S(i,s))}\Big]. \end{split}$$



Figure: Output bound with MMOO arrivals and constant rate server.

Average Improvement is quite decent

$$f_1 = \text{foi}$$

- Single node output with MMOO arrivals and constant rate server
- Randomly chosen parameters in a Monte-Carlo type fashion
- Investigate average and maximum improvement factor:

Standard output bound New output bound

	Improvement Factor
Average	1.66
Maximum	400

Hope for the Delay?

- Direct application to the delay bound is not possible
- In a larger network, multiple output bound computations are necessary to obtain the delay bound
- Need 2^h (h + 1) output bound computations to obtain delay bound on RHS



Application to canonical Topology

• Canonical setting that captures the effect on the delay bound:



Delay Bound can now be improved...



Figure: Delay bound with MMOO arrivals and constant rate server.

...but the Improvement is less significant on average



- As for the output bounds, evaluation by a Monte-Carlo type parameter space exploration
- Compute

Standard delay bound

New delay bound

	Improvement Factor
Average	1.18
Maximum	323



Improvement scales with the Number of Applications



• Lyapunov parameters I_i , i = 1, ..., n-1

Computation Time is scalable



- Investigated computational overhead using Pattern Search for parameter optimization
- Customizable number of parameters

Computation time for $\#$ of flows	2	4	6
new approachstandard apprach	1.99	3.85	6.12

Conclusion

- Improved MGF output bound by inserting Lyapunov's inequality
- New approach is always at least as good as standard bound and is minimally invasive
- Can lead to significant output bound improvement
- Allows for a tighter delay bound in a larger network
- Effect increases with number of output bound computations
- Cost of additional parameter to optimize (but can be conveniently scaled)

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Thank you for your attention! Q & A

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Equivalent Relation

Union Bound

For $X_i \ge 0$:

$$P\left(\max_{i=1,\dots,n} X_{i} > a\right) \stackrel{(UB)}{\leq} \sum_{i=1}^{n} P(X_{i} > a) \stackrel{(Chernoff)}{\leq} e^{-\theta a} \sum_{i=1}^{n} E\left[e^{\theta X_{i}}\right]$$

$$\Leftrightarrow \quad P\left(\max_{i=1,\dots,n} X_{i} > a\right) \stackrel{(Chernoff)}{\leq} e^{-\theta a} E\left[\max_{i=1,\dots,n} e^{\theta X_{i}}\right] \leq e^{-\theta a} \sum_{i=1}^{n} E\left[e^{\theta X_{i}}\right]$$

"quasi-Union bound"